

ON CHARACTERIZATIONS OF LATTICES USING THE GENERALIZED SYMMETRIC BI-DERIVATIONS

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ABSTRACT. In this paper, we introduced the notion of generalized symmetric bi-derivation on lattices and investigated some related properties. We characterized the distributive and modular lattices by generalized symmetric bi-derivations.

1. INTRODUCTION

The lattice algebra plays a significant role in various branches such as information theory, information retrieval, information access controls and cryptanalysis [4, 6, 8, 16]. Recently, the properties of lattices were widely researched. In the theory of rings, the notion of derivation is an important topic to study. After the derivation on a ring was defined by Posner in [15], many researchers studied the derivation theory on various algebraic structures. In [9, 20], authors introduced the notion of derivation on a lattice and discussed some related properties. In [1], Alshehri is introduced the notion of generalized derivation for a lattice and investigated various properties. After the symmetric biderivation on rings was defined in [10, 11] by Maksa, a lot of researchers studied the symmetric biderivation in rings and near-rings [12, 13, 17, 18, 19]. In [7], Çeven applied the notion of symmetric biderivation to lattices and investigated some related properties. The notion of generalized biderivation on rings was introduced by Argaç in [2].

In this paper, we apply the notion of generalized symmetric biderivation to lattices and investigate some related properties which is discussed in [7] and [14]. Also we characterize the distributive and modular lattices by generalized symmetric biderivations.

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2. PRELIMINARIES

Definition 1. ([5]) Let L be a nonempty set endowed with operations " \wedge " and " \vee ". If (L, \wedge, \vee) satisfies the following conditions for all $x, y, z \in L$, then L is called a lattice.

- (a) $x \wedge x = x, x \vee x = x,$
- (b) $x \wedge y = y \wedge x, x \vee y = y \vee x,$
- (c) $(x \wedge y) \wedge z = x \wedge (y \wedge z), (x \vee y) \vee z = x \vee (y \vee z),$
- (d) $(x \wedge y) \vee x = x, (x \vee y) \wedge x = x.$

Definition 2. ([5]) Let (L, \wedge, \vee) be a lattice. A binary relation " \leq " is defined by $x \leq y$ if and only if $x \wedge y = x$ and $x \vee y = y$.

Definition 3. [5] A lattice L is distributive if the following identities hold:

- i) $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z),$
- ii) $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z).$

In any lattice, the above conditions are equivalent.

Definition 4. ([3]) A lattice L is modular if the following identity holds:

$$\text{If } x \leq z, \text{ then } x \vee (y \wedge z) = (x \vee y) \wedge z.$$

Lemma 1 ([20]). *Let (L, \wedge, \vee) be a lattice. Define the binary relation " \leq " as the Definition 2. Then (L, \leq) is a poset and for any $x, y \in L$, $x \wedge y$ is the g.l.b of $\{x, y\}$, and $x \vee y$ is the l.u.b. of $\{x, y\}$.*

Definition 5. Let L be a lattice. A mapping $D(., .) : L \times L \rightarrow L$ is called symmetric if $D(x, y) = D(y, x)$ holds for all $x, y \in L$.

Definition 6. Let L be a lattice. A mapping $d : L \rightarrow L$ defined by $d(x) = D(x, x)$ is called trace of $D(., .)$, where $D(., .) : L \times L \rightarrow L$ is a symmetric mapping.

Definition 7. ([10, 20]) Let L be a lattice and $d : L \rightarrow L$ be a mapping. The mapping d is called a derivation on L , if it satisfies the following condition

$$d(x \wedge y) = (d(x) \wedge y) \vee (x \wedge d(y))$$

for all $x, y, z \in L$.

Definition 8. ([1]) Let L be a lattice. A function $D : L \rightarrow L$ is called a generalized derivation on L if there exists a derivation $d : L \rightarrow L$ such that

$$D(x \wedge y) = (D(x) \wedge y) \vee (x \wedge d(y))$$

for all $x, y, z \in L$.

Definition 9. ([7]) Let L be a lattice and $D : L \times L \rightarrow L$ be a symmetric mapping. We call D a symmetric biderivation on L , if it satisfies the following condition

$$D(x \wedge y, z) = (D(x, z) \wedge y) \vee (x \wedge D(y, z))$$

for all $x, y, z \in L$.

Note that if D is a symmetric biderivation on a lattice L , then the mappings $d_1 : L \rightarrow L, d_1(x) = D(x, y)$ and $d_2 : L \rightarrow L, d_2(y) = D(x, y)$ are derivations on L .

Proposition 1 ([7]). *Let L be a lattice and let d be the trace of symmetric biderivation D . Then the following hold:*

- i) $D(x, y) \leq x$ and $D(x, y) \leq y$,
- ii) $D(x, y) \leq x \wedge y$,
- iii) $d(x) \leq x$,
- iv) $d^2(x) = d(x)$,

for all $x, y \in L$.

3. THE GENERALIZED SYMMETRIC BI-DERIVATIONS ON LATTICES

The following definition introduces the notion of generalized symmetric bi-derivation for a lattice.

Definition 10. Let L be a lattice, $D : L \times L \rightarrow L$ be a symmetric biderivation and $\Delta : L \times L \rightarrow L$ be a symmetric mapping. We call Δ a generalized symmetric biderivation related to D , if it satisfies the following condition

$$\Delta(x \wedge y, z) = (\Delta(x, z) \wedge y) \vee (x \wedge D(y, z))$$

for all $x, y, z \in L$. The mapping $\delta : L \rightarrow L$ defined by $\delta(x) = \Delta(x, x)$ is called the trace of generalized symmetric biderivation Δ .

Obviously, a generalized symmetric biderivation Δ on L satisfies the relation $\Delta(x, y \wedge z) = (\Delta(x, y) \wedge z) \vee (y \wedge D(x, z))$ for all $x, y, z \in L$.

Now we give examples and present some properties for a generalized symmetric biderivation on L .

Example 1. Let L be a lattice with a least element 0. The mapping $D(x, y) = 0$ is a symmetric biderivation on L . Define a mapping on L by $\Delta(x, y) = x \wedge y$ for all $x, y \in L$. Then we can see that Δ is a generalized symmetric biderivation related to D on L .

Example 2. Let L be a lattice with a least element 0 and $a \in L$. The mapping on L defined by $\Delta(x, y) = (x \wedge y) \wedge a$ is a generalized symmetric biderivation related to $D(x, y) = 0$ on L .

Proposition 2. *Let Δ is a generalized symmetric biderivation related to a symmetric biderivation D . Then the mappings $f_1 : L \rightarrow L$, $f_1(x) = \Delta(x, z)$ and $f_2 : L \rightarrow L$, $f_2(y) = \Delta(x, y)$ are generalized derivations on L .*

$$\begin{aligned} f_1(x \wedge y) &= \Delta(x \wedge y, z) \\ \text{Proof. We have} &= (\Delta(x, z) \wedge y) \vee (x \wedge D(y, z)) \\ &= (f_1(x) \wedge y) \vee (x \wedge g_1(y)). \end{aligned}$$

In this equation, the mapping $g_1 : L \rightarrow L$, $g_1(y) = D(y, z)$ is a derivation on L where D is the symmetric biderivation. Hence the mapping f_1 is a generalized derivation on L . \square

Theorem 1. *Let L be a lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ and d be the trace of D . Then*

- (i) $D(x, y) \leq \Delta(x, y)$ for all $x, y \in L$.

If L is distributive lattice, then

- (ii) $\Delta(x, y) \leq x$ and $\Delta(x, y) \leq y$,
- (iii) $\Delta(x, y) \leq x \wedge y$,

- (iv) $d(x) \leq \delta(x) \leq x$
(v) $d(x) = x \implies \delta(x) = x$

for all $x, y \in L$.

Proof. (i) Since

$$\begin{aligned}\Delta(x, y) &= \Delta(x \wedge x, y) \\ &= (\Delta(x, y) \wedge x) \vee (x \wedge D(x, y)) \\ &= (\Delta(x, y) \wedge x) \vee D(x, y), \quad (\text{by Proposition 1 (i)})\end{aligned}$$

it is seen that $D(x, y) \leq \Delta(x, y)$.

(ii) If L is distributive lattice, then we have

$$\begin{aligned}\Delta(x, y) &= (\Delta(x, y) \wedge x) \vee D(x, y) \\ &= (\Delta(x, y) \vee D(x, y)) \wedge (x \vee D(x, y)) \\ &= \Delta(x, y) \wedge x, \quad \text{by (i) and Proposition 1 (i)}.\end{aligned}$$

Hence it is seen that $\Delta(x, y) \leq x$. Since Δ is a symmetric, we have also $\Delta(x, y) \leq y$.

(iii) is clear by (ii).

(iv) is clear by (i) and (ii).

(v) is clear by (iv). □

Corollary 1. *Let L be a distributive lattice and Δ be a generalized symmetric biderivation related to a symmetric biderivation D . Let the least element of L be 0 and the greatest element of L be 1. Then $\Delta(0, x) = \Delta(x, 0) = 0$ and $\Delta(1, x) = \Delta(x, 1) \leq x$ for all $x \in L$.*

Proof. It is trivial from the Theorem 1 (ii). □

Theorem 2. *Let L be a distributive lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ and d be the trace of D . Then*

- (i) $\delta^2(x) = \delta(x)$,
(ii) $\delta(x \wedge y) = (\delta(x) \wedge y) \vee (x \wedge d(y)) \vee D(x, y)$
(iii) $D(x, y) \leq \delta(x \wedge y), \delta(x) \wedge y \leq \delta(x \wedge y), x \wedge d(y) \leq \delta(x \wedge y)$

for all $x, y \in L$.

Proof. (i) Using Proposition 1(i) and Theorem 1 (iv), we have

$$\begin{aligned}\delta^2(x) &= \delta(\delta(x)) \\ &= \delta(x \wedge \delta(x)) \\ &= \Delta(x \wedge \delta(x), x \wedge \delta(x)) \\ &= (\Delta(x, x \wedge \delta(x)) \wedge \delta(x)) \vee (x \wedge D(\delta(x), x \wedge \delta(x))) \\ &= \{[(\Delta(x, x) \wedge \delta(x)) \vee (x \wedge D(\delta(x), x))] \wedge \delta(x)\} \\ &\vee \{x \wedge [(D(x, \delta(x)) \wedge \delta(x)) \vee (x \wedge D(\delta(x), \delta(x)))]\} \\ &= \{[\delta(x) \vee D(\delta(x), x)] \wedge \delta(x)\} \vee \{x \wedge [\delta(x) \vee \delta(x)]\} \\ &= \delta(x).\end{aligned}$$

(ii) Using Proposition 1(i) and Theorem 1 (iv), we have

$$\begin{aligned}\delta(x \wedge y) &= \Delta(x \wedge y, x \wedge y) \\ &= (\Delta(x, x \wedge y) \wedge y) \vee (x \wedge D(y, x \wedge y)) \\ &= \{[(\Delta(x, x) \wedge y) \vee (x \wedge D(x, y))] \wedge y\} \\ &\vee \{x \wedge [(D(x, y) \wedge y) \vee (x \wedge D(y, y))]\} \\ &= \{[(\delta(x) \wedge y) \vee D(x, y)] \wedge y\} \vee \{x \wedge [D(x, y) \vee (x \wedge d(y))]\} \\ &= \{(\delta(x) \wedge y) \vee D(x, y)\} \vee \{D(x, y) \vee (x \wedge d(y))\} \\ &= (\delta(x) \wedge y) \vee (x \wedge d(y)) \vee D(x, y)\end{aligned}$$

(iii) It is clear by (ii). \square

Theorem 3. *Let L be a distributive lattice and Δ be a generalized symmetric biderivation related to a symmetric biderivation D and δ be the trace of Δ . Then*

$$d(x) \wedge d(y) \leq \delta(x) \wedge \delta(y) \leq \delta(x \wedge y)$$

for all $x, y \in L$.

Proof. By Theorem 2 (iii) and Theorem 1 (iv), since $\delta(x) \wedge y \leq \delta(x \wedge y)$ and $\delta(y) \leq y$, we get $\delta(x) \wedge \delta(y) \leq \delta(x) \wedge y \leq \delta(x \wedge y)$. Using Theorem 1 (iv), we have $d(x) \wedge d(y) \leq \delta(x) \wedge \delta(y)$. \square

Theorem 4. *Let L be a distributive lattice and Δ be a generalized symmetric biderivation related to a symmetric biderivation D . Let the least element of L be 0 and the greatest element of L be 1, then*

- (i) if $x \leq \delta(1)$, then $\delta(x) = x$,
- (ii) if $x \geq \delta(1)$, then $\delta(1) \leq \delta(x)$,
- (iii) if $x \leq y$ and $d(y) = y$, then $\delta(x) = x$.

Proof. From Theorem 2 (iii), we have $\delta(1) \wedge x \leq \delta(x)$. Hence

- (i) if $x \leq \delta(1)$, then we get $x \leq \delta(x)$. Using Theorem 1 (iv), we get $\delta(x) = x$.
- (ii) If $x \geq \delta(1)$, then we have $\delta(1) \leq \delta(x)$.
- (iii) Since $\delta(x) \leq x \leq y$, $d(y) = y$, $D(x, y) \leq x$ we have, by Theorem 2 (ii),

$$\begin{aligned} \delta(x) &= \delta(x \wedge y) \\ &= (\delta(x) \wedge y) \vee (x \wedge d(y)) \vee D(x, y) \\ &= \delta(x) \vee x \vee D(x, y) \\ &= x. \end{aligned}$$

In every lattice L with the least element 0, the mapping $D(x, y) = 0$ is a symmetric biderivation. The mapping $\Delta(x, y) = x \vee y$ in any lattice L is not a generalized symmetric biderivation related to $D(x, y) = 0$. Then we have the following Corollary:

Corollary 2. *In a lattice with the least element 0, if the mapping $\Delta(x, y) = x \vee y$ related to $D(x, y) = 0$ is a generalized symmetric biderivation, then the lattice L is modular lattice.*

Proof. From the equality $\Delta(y \wedge z, x) = (\Delta(y, x) \wedge z) \vee (y \wedge D(z, x))$, we have $x \vee (y \wedge z) = (x \vee y) \wedge z$ for all $x, y, z \in L$, hence L is a modular lattice. \square

Definition 11. Let L be a lattice. The mapping Δ satisfying $\Delta(x \vee y, z) = \Delta(x, z) \vee \Delta(y, z)$ for all $x, y \in L$, is called a jointive mapping.

Theorem 5. *Let L be a lattice and Δ be a jointive and symmetric mapping with the trace δ on L . Then*

- (i) $\delta(x \vee y) = \delta(x) \vee \delta(y) \vee \Delta(x, y)$,
- (ii) $\delta(x) \vee \delta(y) \leq \delta(x \vee y)$

for all $x, y \in L$.

Proof. (i) By the definition of jointive mapping and symmetry, we have

$$\begin{aligned} \delta(x \vee y) &= \Delta(x \vee y, x \vee y) \\ &= \Delta(x, x) \vee \Delta(x, y) \vee \Delta(y, y) \\ &= \delta(x) \vee \delta(y) \vee \Delta(x, y). \end{aligned}$$

(ii) it is clear from (i). \square

Proposition 3. *If the mapping $\Delta(x, y) = x \wedge y$ in a lattice L is also a joinitive mapping, then the lattice L is a distributive lattice.*

Proof. Using the equality $\Delta(x \vee y, z) = \Delta(x, z) \vee \Delta(y, z)$, we have $(x \vee y) \wedge z = (x \wedge z) \vee (y \wedge z)$ for all $x, y \in L$. So L is a distributive lattice. \square

Let L be a distributive lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ and d be the trace of D . Denote $Fix_d(L) = \{x \in L : d(x) = x\}$. By Theorem 1 (v), $x \in Fix_d(L)$ implies that $x \in Fix_\delta(L)$ for all $x \in L$. That is, $Fix_d(L) \subseteq Fix_\delta(L)$. Furthermore, from the Theorem 4 (iii), $x \in Fix_d(L)$ and $y \leq x$ implies that $y \in Fix_\delta(L)$.

Definition 12. Let L be a lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ . If $x \leq y$ implies $\delta(x) \leq \delta(y)$, then δ is called an isotone mapping.

Proposition 4. *Let L be a distributive lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ . If δ is isotone and $x, y \in Fix_\delta(L)$, then $\delta(x \vee y) = x \vee y$ for all $x, y \in L$.*

Proof. Since $x \leq x \vee y$ and $y \leq x \vee y$ and δ is isotone, we have $\delta(x) \leq \delta(x \vee y)$ and $\delta(y) \leq \delta(x \vee y)$. So it is seen that $\delta(x) \vee \delta(y) \leq \delta(x \vee y)$ and since $x, y \in Fix_\delta(L)$, $x \vee y \leq \delta(x \vee y)$. By Theorem 1 (iv), since $\delta(x \vee y) \leq x \vee y$, we obtain $\delta(x \vee y) = x \vee y$. \square

Proposition 5. *Let L be a distributive lattice, Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ and d be the trace of D . Then $1 \in Fix_\delta(L)$ if and only if δ is an identity mapping.*

Proof. If $1 \in Fix_\delta(L)$, since $\delta(1) = 1$, by Theorem 2, we have

$$\delta(x) = \delta(1 \wedge x) = (\delta(1) \wedge x) \vee (1 \wedge d(x)) \vee D(1, x) = x \vee d(x) \vee D(1, x) = x.$$

Converse is trivial. \square

Theorem 6. *Let L be a distributive lattice, Δ_1 and Δ_2 be generalized symmetric biderivations related to a same symmetric biderivation D . The mapping $\Delta_1 \wedge \Delta_2$ defined by $(\Delta_1 \wedge \Delta_2)(x, y) = \Delta_1(x, y) \wedge \Delta_2(x, y)$, is a generalized symmetric biderivation related to the symmetric biderivation D .*

Proof. Since

$$\begin{aligned} (\Delta_1 \wedge \Delta_2)(x \wedge y, z) &= \Delta_1(x \wedge y, z) \wedge \Delta_2(x \wedge y, z) \\ &= [(\Delta_1(x, z) \wedge y) \vee (x \wedge D(y, z))] \\ &\quad \wedge [(\Delta_2(x, z) \wedge y) \vee (x \wedge D(y, z))] \\ &= [(\Delta_1(x, z) \wedge y) \wedge (\Delta_2(x, z) \wedge y)] \vee (x \wedge D(y, z)) \\ &= (\Delta_1(x, z) \wedge \Delta_2(x, z) \wedge y) \vee (x \wedge D(y, z)) \\ &= ((\Delta_1 \wedge \Delta_2)(x, z) \wedge y) \vee (x \wedge D(y, z)), \end{aligned}$$

so the Theorem is true. \square

Theorem 7. *Let L be a distributive lattice, Δ_1 and Δ_2 be generalized symmetric biderivations related to a same symmetric biderivation D . The mapping $\Delta_1 \vee \Delta_2$ defined by $(\Delta_1 \vee \Delta_2)(x, y) = \Delta_1(x, y) \vee \Delta_2(x, y)$, is a generalized symmetric biderivation related to the symmetric biderivation D .*

Proof. Since

$$\begin{aligned}
(\Delta_1 \vee \Delta_2)(x \wedge y, z) &= \Delta_1(x \wedge y, z) \vee \Delta_2(x \wedge y, z) \\
&= (\Delta_1(x, z) \wedge y) \vee (x \wedge D(y, z)) \\
&\quad \vee (\Delta_2(x, z) \wedge y) \vee (x \wedge D(y, z)) \\
&= (\Delta_1(x, z) \wedge y) \vee (\Delta_2(x, z) \wedge y) \vee (x \wedge D(y, z)) \\
&= ((\Delta_1(x, z) \vee \Delta_2(x, z)) \wedge y) \vee (x \wedge D(y, z)) \\
&= ((\Delta_1 \vee \Delta_2)(x, z) \wedge y) \vee (x \wedge D(y, z)),
\end{aligned}$$

so the Theorem is true. \square

Proposition 6. *Let L be a distributive lattice, Δ_1 and Δ_2 be generalized symmetric biderivations, δ_1 be the trace of Δ_1 and δ_2 be the trace of Δ_2 . If δ_1 and δ_2 are isotone mapping, then $\delta_1 = \delta_2$ if and only if $Fix_{\delta_1}(L) = Fix_{\delta_2}(L)$.*

Proof. Let $Fix_{\delta_1}(L) = Fix_{\delta_2}(L)$. If $x \in Fix_{\delta_1}(L)$, since $\delta_1(\delta_1(x)) = \delta_1(x)$, we have $\delta_1(x) \in Fix_{\delta_1}(L) = Fix_{\delta_2}(L)$. Hence $\delta_2(\delta_1(x)) = \delta_1(x)$. Similarly, we see that $\delta_1(\delta_2(x)) = \delta_2(x)$. Since δ_1 and δ_2 are isotone mapping and $\delta_1(x) \leq x, \delta_2(x) \leq x$, we get $\delta_2(\delta_1(x)) \leq \delta_2(x) = \delta_1(\delta_2(x))$ and $\delta_1(\delta_2(x)) \leq \delta_1(x) = \delta_2(\delta_1(x))$. So $\delta_1(\delta_2(x)) = \delta_2(\delta_1(x))$. Therefore we obtain $\delta_1(x) = \delta_2(\delta_1(x)) = \delta_1(\delta_2(x)) = \delta_2(x)$, that is, $\delta_1 = \delta_2$. The converse is trivial. \square

Proposition 7. *Let L be a distributive lattice and Δ be a generalized symmetric biderivation related to a symmetric biderivation D , δ be the trace of Δ and the greatest element of L be 1. Then the following conditions are equivalent:*

- (i) δ is an isotone mapping
- (ii) $\delta(x) \vee \delta(y) \leq \delta(x \vee y)$
- (iii) $\delta(x \wedge y) = \delta(x) \wedge \delta(y)$
- (iv) $\delta(x) = x \wedge \delta(1)$ for all $x, y \in L$.

Proof. (i) \implies (ii): Since $x \leq x \vee y$ and $y \leq x \vee y$ and δ is an isotone mapping, we have $\delta(x) \leq \delta(x \vee y)$ and $\delta(y) \leq \delta(x \vee y)$, so $\delta(x) \vee \delta(y) \leq \delta(x \vee y)$.

(ii) \implies (i): Let $\delta(x) \vee \delta(y) \leq \delta(x \vee y)$ and $x \leq y$. Since $x \vee y = y$, we have $\delta(x) \vee \delta(y) \leq \delta(y)$. Also it is known that $\delta(y) \leq \delta(x) \vee \delta(y)$. Hence we obtain $\delta(x) \vee \delta(y) = \delta(y)$, so $\delta(x) \leq \delta(y)$.

(i) \implies (iii): Since $x \wedge y \leq x$ and $x \wedge y \leq y$ and δ is an isotone mapping, we have $\delta(x \wedge y) \leq \delta(x)$ and $\delta(x \wedge y) \leq \delta(y)$ and so $\delta(x \wedge y) \leq \delta(x) \wedge \delta(y)$. By Theorem 3 (ii), we have $\delta(x) \wedge \delta(y) \leq \delta(x \wedge y)$. Hence $\delta(x) \wedge \delta(y) = \delta(x \wedge y)$.

(iii) \implies (i): Let $\delta(x \wedge y) = \delta(x) \wedge \delta(y)$ and $x \leq y$. Since $x \wedge y = x$, we get $\delta(x) = \delta(x \wedge y) = \delta(x) \wedge \delta(y) \leq \delta(y)$.

(i) \implies (iv): Since $x \leq 1$ and δ is an isotone mapping, we have $\delta(x) \leq \delta(1)$. By Theorem 1 (iv), since $\delta(x) \leq x$, we get $\delta(x) \leq x \wedge \delta(1)$. By Theorem 2 (ii), we have $\delta(x) \wedge y \leq \delta(x \wedge y)$. Taking $x = 1$, we get $\delta(1) \wedge y \leq \delta(y)$ for all $y \in L$. Hence we have $\delta(x) = x \wedge \delta(1)$.

(iv) \implies (iv): Let $\delta(x) = x \wedge \delta(1)$ and $x \leq y$. Since $x \wedge y = x$, we have $\delta(x) = \delta(x \wedge y) = (x \wedge y) \wedge 1 = (x \wedge 1) \wedge (y \wedge 1) = \delta(x) \wedge \delta(y)$. Hence $\delta(x) \leq \delta(y)$. \square

Conflict of Interests

The author declares that there is no conflict of interest regarding the publication of this paper.

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